

Various feedback on the course:

By Ludovic Schwartz :

a convex function  $f: X \rightarrow \mathbb{R}$  (or even  $\mathbb{C}_1$ ) does not necessarily achieve its infimum,

hence we safely and rightfully considered  $p_f^*$  as an approximate  $\arg\inf$  in the last proof of Lecture 1.

Ex:  $f(x) = x^2 + \frac{1}{x} \text{ for } x > 0$ , infimum equal to 0, never achieved

(take  $f(x) = x^2 + \frac{1}{x} \text{ for } x < 0$  for an exp-concave example)

By Hassan Saber :

in the last proof it is easy to directly optimize over  $\varepsilon \in (0,1)$  the obtained bound:

$$B(\varepsilon) = T \ln(1-\varepsilon) + (N-1) \ln \varepsilon$$

$$B'(\varepsilon) = \frac{-T}{1-\varepsilon} + \frac{N-1}{\varepsilon}$$

$B''(\varepsilon) < 0$  thus the  $\varepsilon^*$  s.t.  $B'(\varepsilon^*) = 0$  is the global minimum of  $B$

$$\text{We have } B'(\varepsilon^*) = 0 \Leftrightarrow \frac{\varepsilon^*}{N-1} = \frac{1}{T} - \frac{\varepsilon^*}{T} \Leftrightarrow \varepsilon^* = \frac{1/T}{1/T + 1/(N-1)} = \frac{N-1}{T+N-1}$$

We thus get the sharper final bound

$$\begin{aligned} \frac{1}{T} \left( T \ln \left( \frac{T+N-1}{T} \right) + (N-1) \ln \left( \frac{T+N-1}{N-1} \right) \right) &= \frac{T+N-1}{T} H \left( \frac{N-1}{T+N-1} \right) \\ &= \frac{1}{T} \left( (N-1) \underbrace{\ln \left( 1 + \frac{N-1}{T} \right)}_{\leq (N-1)^2/T} + (N-1) \ln \left( 1 + \frac{T}{N-1} \right) \right) \end{aligned}$$

where  $H(x) = x \ln x + (1-x) \ln(1-x)$   
is the binary entropy

\* However \* both bounds are  $\sim \frac{N-1}{T} \ln T$  as  $T \rightarrow +\infty$ .

Exercise #1 / Various considerations around the notion of regret.

1) ( $N=2$  is enough)

Time  $t=1 \quad t=2 \quad t=3 \quad t=4 \quad t=5 \quad t=6 \quad \dots$

$$l_{1t} \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0$$

$$\sum_{s=1}^{t-1} l_{1s} \quad 1 \quad 1 \quad 2 \quad 2 \quad 3$$

$$l_{2t} \quad \frac{1}{2} \quad 1 \quad 0 \quad 1 \quad 0 \quad 1$$

$$\sum_{s=1}^{t-1} l_{2s} \quad \frac{1}{2} \quad 1+\frac{1}{2} \quad 1+\frac{1}{2} \quad 2+\frac{1}{2} \quad 2+\frac{1}{2}$$

$$\text{Leader over } s=1, \dots, t-1 \quad 2 \quad 1 \quad 2 \quad 1 \quad 2$$

$$p_{1t} \quad \frac{1}{2} \quad 0 \quad 1 \quad 0 \quad 1 \quad 0$$

$$p_{2t} \quad \frac{1}{2} \quad 1 \quad 0 \quad 1 \quad 0 \quad 1$$

$$\sum_{j=1}^2 p_j l_{jt} \quad \frac{3}{4} \quad 1 \quad 1 \quad 1 \quad 1 \quad 1$$

$$\sum_{s=1}^t \sum_{j=1}^2 p_j l_{js} \quad \frac{3}{4} \quad 1+\frac{3}{4} \quad 2+\frac{3}{4} \quad 3+\frac{3}{4} \quad 4+\frac{3}{4} \quad 5+\frac{3}{4}$$

Follow  
Be leader fails  
is a strategy  
ENT is a smoothed alternative.  
is a minima  
↑

We have

$$\sum_{t=1}^T \sum_{j=1}^2 p_j l_{jt} = T - 1 + \frac{3}{4}$$

$$\sum_{t=1}^T l_{1t} = \begin{cases} T/2 & \text{if } T \text{ even} \\ \frac{T-1}{2} & \text{if } T \text{ odd} \end{cases} \quad (\text{lower integer part})$$

$$\sum_{t=1}^T l_{2t} = \begin{cases} (T-1)/2 + \frac{1}{2} & \text{if } T \text{ odd} \\ T/2 - 1 + \frac{1}{2} & \text{if } T \text{ even} \end{cases}$$

$$\min \left\{ \sum_{t=1}^T l_{1t}, \sum_{t=1}^T l_{2t} \right\} \leq \frac{T}{2}$$

$$R_T \geq T - 1 + \frac{3}{4} - \frac{T}{2} = \frac{T}{2} - \frac{1}{4} \neq o(T)$$

2) It suffices to consider only the sequences in  $\{q_1\}_{t=1}^N$ :

For any given strategy, we denote for  $t \geq 2$ :

$$k_t^* \in \operatorname{argmax}_{k \in \{1..N\}} p_{kt} \rightarrow \begin{cases} \text{in particular,} \\ p_{k_t^*, t} \geq \frac{1}{N} \end{cases}$$

Note that  $k_t^*$  depends only on  $l_{js}, j \in \{1..N\}, s \in \{1..t-1\}$

The sequence

$$\begin{cases} l_{k_t^*, t} = 1 \\ l_{jt} = 0 \quad j \neq k_t^* \end{cases}$$

is such that:

$$\sum_{t=1}^T \sum_j p_{jt} l_{jt} \geq \underbrace{\sum_{t=1}^T p_{k_t^*, t} l_{k_t^*, t}}_{\geq \frac{1}{N}} \geq \frac{T}{N}$$

while

$$\sum_{t=1}^T \min_j l_{jt} = 0.$$

Thus, any strategy can be such that

$$\sup_{\substack{(l_{kt} - l_{nt}) \in \{0, 1\}^N \\ t=1..T}} \left\{ \sum_{t=1}^T p_{jt} l_{jt} - \sum_t \min_k l_{kt} \right\} = o(T)$$

3) In the proof, instead of applying Hoeffding's lemma

$$\ln E[e^{\eta X}] \leq \eta E[X] + \frac{\eta^2}{8} (M-m)^2$$

we apply Jensen's inequality:

$$\ln E[e^{\eta X}] \geq \eta E[X]$$

(valid  $\forall \eta \in \mathbb{R}$  and all variables  $X$   
s.t.  $X$  is integrable)

Then  $\sum_j p_j f_j \ln f_j = -\frac{1}{\eta} \underbrace{\left( -\eta \sum_j p_j f_j \right)}_{< \ln \sum_j p_j e^{-\eta f_j}} \geq -\frac{1}{\eta} \ln \sum_j p_j e^{-\eta f_j}$

Now, with the same telescoping argument:

$$\begin{aligned} \sum_{t=1}^T \sum_{j=1}^N p_j f_j \ln f_j &\geq -\frac{1}{\eta} \ln \frac{\sum_{j=1}^N e^{-\eta \sum_{t=1}^T f_j t}}{N} \\ &\geq -\frac{1}{\eta} \ln \max_{j=1 \dots N} e^{-\eta \sum_{t=1}^T f_j t} \\ &= \min_{j=1 \dots N} \sum_{t=1}^T f_j t \end{aligned}$$

(upper bound by the average)

That,  $\forall \eta > 0$ ,  $R_T = \sum_{t,j} p_j f_j \ln f_j - \min_k \sum_{t=1}^T f_k t \geq 0$

Exercise #3Convex loss functions and comparison to the best convex vector

Strategy at hand:

$$\eta > 0 \quad \text{and}$$

$$p_t = \int_X p e^{-\eta \sum_{s=1}^{t-1} l_s(p)} d\mu(p) / \int_X e^{-\eta \sum_{s=1}^{t-1} l_s(p)} d\mu(p)$$

$$= \int_X p d\mu_t(p) \quad \text{where } \frac{d\mu_t}{dp}(p) = \frac{e^{-\eta \sum_{s=1}^{t-1} l_s(p)}}{\int_X e^{-\eta \sum_{s=1}^{t-1} l_s(q)} d\mu(q)}$$

$$1) \quad l_t(p_t) = l_t\left(\int p d\mu_t(p)\right) \stackrel{\text{Jensen}}{\leq} \int_X l_t(p) d\mu_t(p)$$

$\stackrel{\text{as for EWA}}{\leq} -\frac{1}{\eta} \ln \underbrace{\int e^{-\eta l_t(p)} d\mu_t(p)}_{\ln \frac{\int e^{-\eta \sum_{s=1}^t l_s(p)} d\mu(p)}{\int e^{-\eta \sum_{s=1}^{t-1} l_s(p)} d\mu(p)}} + \frac{(M-m)^2}{8\eta}$

Summing over  $t=1, \dots, T$ , a telescoping sum appears:

$$\sum_{t=1}^T l_t(p_t) \leq -\frac{1}{\eta} \ln \underbrace{\frac{\int e^{-\eta \sum_{t=1}^T l_t(p)} d\mu(p)}{1}}_{\text{can be bounded using the same techniques as for exp-concave losses, but the proof needs to be slightly adapted:}} + \frac{(M-m)^2}{8\eta} \eta T$$

 $\delta > 0$  and  $p_S^*$  s.t.

$$\inf_{p \in X} \sum_{t=1}^T l_t(p) \leq \delta + \sum_{t=1}^T l_t(p_S^*)$$

$$\varepsilon > 0 \quad \text{and} \quad \Delta_{S,\varepsilon}^* = \{ (1-\varepsilon)p_S^* + \varepsilon r, \quad r \in X \}$$

$$\text{We still have } \mu(\Delta_{S,\varepsilon}^*) = \varepsilon^{N-1}$$

But for  $p = (1-\varepsilon) p_S^* + \varepsilon r$  we can only resort to convexity:

$$\begin{aligned} l_t(p) &\leq (1-\varepsilon) l_t(p_S^*) + \varepsilon l_t(r) \\ &\leq l_t(p_S^*) + \varepsilon \underbrace{(l_t(r) - l_t(p_S^*))}_{\leq M-m \text{ since } l_t \text{ takes values in } [m, M]} \end{aligned}$$

$$e^{-\eta l_t(p)} \geq e^{-\eta l_t(p_S^*)} e^{-\eta \varepsilon (M-m)}$$

$l_t$  takes values in  $[m, M]$   
by assumption

Putting all things together:

$$\int_X e^{-\eta \sum_{t=1}^T l_t(p)} d\mu(p) \geq e^{-\eta \sum_{t=1}^T l_t(p_S^*)} \times e^{-\eta \varepsilon (M-m)T} \times \varepsilon^{N-1}$$

↑  
 integral  
 only over  
 $\Delta_S^*$

Substituting above and taking  $\inf_{\varepsilon}$ :

$$\begin{aligned} \sum_{t=1}^T l_t(p_t) &\leq \underbrace{\sum_{t=1}^T l_t(p_S^*)}_{\leq \delta + \inf_p \sum_{t=1}^T l_t(p)} + \inf_{\varepsilon \in (0, 1)} \left\{ \varepsilon(M-m)T - \frac{N-1}{\eta} \ln \varepsilon \right\} \\ &\quad + \frac{(M-m)^2}{8} \eta T \end{aligned}$$

We let  $\delta \downarrow 0$  to conclude:

$$\begin{aligned} \sum_{t=1}^T l_t(p_t) - \inf_{p \in X} \sum_{t=1}^T l_t(p) &\leq \frac{(M-m)^2}{8} \eta T + \inf_{\varepsilon \in (0, 1)} \left( \varepsilon(M-m)T - \frac{N-1}{\eta} \ln \varepsilon \right) \end{aligned}$$

2) Optimize first over  $\varepsilon$ :

$g$  strictly convex,  
a unique minimizer

on  $(0, +\infty)$  at  $\varepsilon$  s.t.

$$\left\{ \begin{array}{l} g(\varepsilon) = \varepsilon(M-m)T - \frac{N-1}{\eta} \ln \varepsilon \\ g'(\varepsilon) = (M-m)T - \frac{N-1}{\eta \varepsilon} \\ g''(\varepsilon) = \frac{N-1}{\eta \varepsilon^2} > 0 \end{array} \right.$$

$$g'(\varepsilon) = 0 \Leftrightarrow \varepsilon = \frac{N-1}{\eta(M-m)T}$$

⚠ but question is whether this  $\varepsilon$  is in  $(0, 1)$ !

I tried with this value of  $\varepsilon$  (which is ok for large  $T$ ) but couldn't get a simple and readable  $O(\sqrt{NT \ln T})$  bound.

Let's not optimize over  $\varepsilon$  and take an arbitrary choice:  $\varepsilon = 1/\sqrt{T}$

$$\text{The bound is } \leq \frac{(M-m)^2}{8} \eta T + (M-m)\sqrt{T} + \frac{N-1}{2\eta} \ln T$$

$$\text{Optimal value for } \eta : \quad \eta^* \text{ s.t.} \quad \left( \frac{(M-m)^2}{8} T \right) \eta^* = \frac{N-1}{2\eta^*} \ln T$$

(as seen  
in class)

and for this  $\eta^*$ , the sum is  $2 \times \sqrt{\text{the product}}$

$$= \underbrace{\frac{2}{\sqrt{16}}}_{= \frac{1}{2}} (M-m) \sqrt{(N-1)T \ln T}$$

$$\text{Final bound : } \frac{1}{2} (M-m) \sqrt{(N-1)T \ln T} + (M-m) \sqrt{T}.$$

→ If you can proceed better, please send me your solution (and you may be rewarded with bonus points at the exam).