

Let's first complete the proof of the Remark:

[“Hoeffding-Azuma inequality with a random number of summands”]

Setting: Probability distributions  $\nu_1, \dots, \nu_k$  over  $[0,1]$   
with respective expectations  $\mu_1, \dots, \mu_k$

At each round,  $I_t \in \{1, \dots, k\}$  is picked in a  $\sigma(Y_1, \dots, Y_{t-1})$ -measurable way

then  $Y_t$  is drawn independently at random according to  $\nu_{I_t}$ , given  $I_t$

$$\text{i.e.: } Y_t | I_t \sim \nu_{I_t}$$

We denote  $N_a(t) = \sum_{s=1}^t \mathbb{1}_{\{I_s=a\}}$  and assume that each arm  $a$  was pulled once in the first  $K$  rounds, so that:

$$N_a(t) \geq 1 \quad \forall t \geq K$$

Then, for  $t \geq K$ :

$$\hat{\mu}_{a,t} = \frac{1}{N_a(t)} \sum_{s=1}^t Y_s \mathbb{1}_{\{I_s=a\}}$$

Lemma:  $\forall \delta \in (0,1)$ ,  $P\left\{ \mu_a > \hat{\mu}_{a,t} + \sqrt{\frac{\ln(1/\delta)}{2N_a(t)}} \right\} \geq 1 - \delta$

(and a similar symmetric statement with  $\mu_a < \hat{\mu}_{a,t} + \sqrt{\cdot}$ )

The proof will be based on the fact that  $(Z_t)_{t \geq 0}$ , where

$$Z_t = \sum_{s=1}^t (Y_s - \mu_a) \mathbb{1}_{\{I_s=a\}}$$

is a martingale w.r.t.  $(\mathcal{F}_t) = (\sigma(Y_1, \dots, Y_t))_{t \geq 0}$ , which we already proved last time:

$$\begin{aligned} \mathbb{E}\left[ (Y_t - \mu_a) \mathbb{1}_{\{I_t=a\}} \mid Y_1, \dots, Y_{t-1} \right] &= \mathbb{E}\left[ (Y_t - \mu_a) \mathbb{1}_{\{I_t=a\}} \mid I_t, Y_1, \dots, Y_{t-1} \right] \\ &= (\mathbb{E}[Y_t \mid I_t, Y_1, \dots, Y_{t-1}] - \mu_a) \mathbb{1}_{\{I_t=a\}} \\ &\quad \text{where we used the bandit model} \\ &= (\mu_{I_t} - \mu_a) \mathbb{1}_{\{I_t=a\}} = 0 \text{ a.s.} \end{aligned}$$

Remark: How does this bound compare to what the classical version of the Hoeffding-Azuma says?

Martingale increment  $(Y_t - \mu_a) \mathbb{1}_{\{I_t=a\}}$  bounded between

$$a_t = -\mu_a \text{ and } b_t = 1 - \mu_a$$

so that

$$(\text{actually in the version I stated, I can take } \leq \text{ or } \geq) \quad (b_t - a_t)^2 = 1$$

$$1 - ts \leq \mathbb{P}\left\{ Z_t < \sqrt{\frac{t}{2} \ln \frac{1}{ts}} \right\} = \mathbb{P}\left\{ N_a(t) (\hat{\mu}_a - \mu_a) < \sqrt{\frac{t}{2} \ln \frac{1}{ts}} \right\}$$

$$= \mathbb{P}\left\{ \hat{\mu}_a - \frac{\sqrt{\frac{t}{2} \ln(1/ts)}}{\sqrt{N_a(t)}} < \mu_a \right\}$$

versus the bound of our lemma:  $1 - ts \leq \mathbb{P}\left\{ \hat{\mu}_a - \sqrt{\frac{\ln(1/s)}{2N_a(t)}} < \mu_a \right\}$

The proposed deviations essentially differ from a  $\sqrt{\frac{t}{N_a(t)}}$  factor, and it is so nice to get rid of it!

Proof: (1) We prove that  $\forall x \in \mathbb{R}$ ,  $\mathbb{E}[e^{xZ_t - x^2/8 N_a(t)}] \leq 1$

We do so by showing that  $M_t = \exp(xZ_t - \frac{x^2}{8} N_a(t))$  is a supermartingale, so that  $\mathbb{E}[M_t] \leq \mathbb{E}[M_0] = 1$ .

Indeed, by the conditional version of Hoeffding's lemma,

$$\mathbb{E}[e^{x(Y_t - \mu_a) \mathbf{1}_{\{I_t=a\}}} | \mathcal{F}_{t-1}] \leq e^{x^2/8} \quad \text{a.s.} \quad \begin{matrix} \text{but we} \\ \text{can do better!} \end{matrix}$$

Since  $I_t$  and thus also  $\mathbf{1}_{\{I_t=a\}}$  are  $\mathcal{F}_{t-1}$ -measurable, we get:

$$\begin{aligned} \mathbb{E}[e^{x(Y_t - \mu_a) \mathbf{1}_{\{I_t=a\}}} | \mathcal{F}_{t-1}] &= \mathbb{E}[e^{x(Y_t - \mu_a) \mathbf{1}_{\{I_t=a\}} (\mathbf{1}_{\{I_t=a\}} + \mathbf{1}_{\{I_t \neq a\}})} | \mathcal{F}_{t-1}] \\ &= \mathbb{E}[e^{x(Y_t - \mu_a) \mathbf{1}_{\{I_t=a\}}} | \mathcal{F}_{t-1}] \mathbf{1}_{\{I_t=a\}} + e^0 \mathbf{1}_{\{I_t \neq a\}} \\ &\stackrel{\text{given what we had before}}{\leq} e^{x^2/8} \mathbf{1}_{\{I_t=a\}} + \mathbf{1}_{\{I_t \neq a\}} = \exp\left(\frac{x^2}{8} \mathbf{1}_{\{I_t=a\}}\right) \end{aligned}$$

Put differently,  $\mathbb{E}[e^{x(Y_t - \mu_a) \mathbf{1}_{\{I_t=a\}}} - \frac{x^2}{8} \mathbf{1}_{\{I_t=a\}} | \mathcal{F}_{t-1}] \leq 1$

which entails that

$$\begin{aligned} &\exp\left(x \sum_{s=1}^t (Y_s - \mu_a) \mathbf{1}_{\{I_s=a\}} - \frac{x^2}{8} \sum_{s=1}^t \mathbf{1}_{\{I_s=a\}}\right) \\ &= \exp(xZ_t - \frac{x^2}{8} N_a(t)) = M_t \end{aligned}$$

is a supermartingale wrt  $\mathcal{F}_t = \sigma(Y_1, \dots, Y_t)$ .

(2) We prove that  $\forall \varepsilon > 0, \forall l \geq 1, \mathbb{P}\{Z_t \geq \varepsilon \text{ and } N_a(t) = l\} \leq \exp(-2\varepsilon^2/l)$

Indeed, by a Markov-Chernoff bounding,

$\forall x > 0,$

$$\begin{aligned} \mathbb{P}\{Z_t \geq \varepsilon \text{ and } N_a(t) = l\} &\leq e^{-x\varepsilon} \mathbb{E}[e^{xZ_t} \mathbb{1}_{\{N_a(t) = l\}}] \\ &= e^{-x\varepsilon + \frac{x^2l}{8}} \mathbb{E}[e^{xZ_t - \frac{x^2}{8}N_a(t)} \mathbb{1}_{\{N_a(t) = l\}}] \\ &\leq e^{-x\varepsilon + \frac{x^2l}{8}} \underbrace{\mathbb{E}[e^{xZ_t - \frac{x^2}{8}N_a(t)}]}_{\leq 1 \text{ by (1)}} \end{aligned}$$

Optimizing over  $x > 0$

(take  $x = 4\varepsilon/l$ ) yields the claimed bound.

(3) Conclusion: we prove that  $\mathbb{P}\{\mu_a \leq \hat{f}_{ab} - \sqrt{\frac{\ln(1/\delta)}{2N_a(t)}}\} \leq \delta$

Indeed, by distinguishing according to the values taken by  $N_a(t)$ :

$$\begin{aligned} &\mathbb{P}\{\mu_a \leq \hat{f}_{ab} - \sqrt{\frac{\ln(1/\delta)}{2N_a(t)}}\} \\ &= \sum_{l=1}^t \mathbb{P}\{N_a(t) = l \text{ and } \mu_a \leq \hat{f}_{ab} - \sqrt{\frac{\ln(1/\delta)}{2l}}\} \\ &= \sum_{l=1}^t \mathbb{P}\{N_a(t) = l \text{ and } \frac{Z_t}{N_a(t)} \geq \sqrt{\frac{\ln(1/\delta)}{2l}}\} \\ &= \sum_{l=1}^t \mathbb{P}\{N_a(t) = l \text{ and } Z_t \geq \sqrt{l \ln(1/\delta)/2}\} \\ &\stackrel{\text{by (2)}}{\leq} \sum_{l=1}^t \exp(-2(l \ln(1/\delta)/2)/l) = t\delta. \end{aligned}$$

(  $\sum_{l=1}^{t-K+1}$  would be enough )

Stochastic bandits :What about arms indexed by a continuum?Setting 1 :Arms indexed by  $x \in A$ , where  $A$  is some possibly large setWith each arm  $x \in A$  is associated a probability distribution $\tilde{\nu}_x$  over  $\mathbb{R}$  s.t.  $E(\tilde{\nu}_x)$  existsAt each round, the decision-maker picks  $I_t \in A$ ,gets a reward  $y_t$  drawn at random according to  $\tilde{\nu}_{I_t}$ (given  $I_t$ ); and this is the only feedback she gets.Definition:  $f: x \in A \mapsto E(\tilde{\nu}_x)$  is the mean-payoff function

Regret :

$$\bar{R}_T = T \sup_{x \in A} f(x) - E\left[\sum_{t=1}^T y_t\right]$$

Setting 2 : [special case]  $\rightarrow$  Noisy optimization of a function.We fix  $f: A \rightarrow \mathbb{R}$ 

The noise is given by a sequence of iid random variables

 $\varepsilon_1, \varepsilon_2, \dots$ When  $I_t \in A$  is picked,  $y_t = f(I_t) + \varepsilon_t$ 

$\hookrightarrow$  Special case of setting #1 where  $\tilde{\nu}_x$  is the distribution of  $f(x) + \varepsilon_1$  (all these distributions have the same shape given by the common distribution of the  $\varepsilon_j$ )

We of course need conditions for the regret to be minimized.

Definition: Let  $\mathcal{F}$  be a set of possible bandit problems  $\mathcal{F} = (\tilde{\nu}_x)_{x \in A}$  $\rightarrow$  The regret can be controlled (in a non-uniform way) against  $\mathcal{F}$  if:

we also say that

 $(A, \mathcal{F})$  is tractablethere exists a strategy s.t.  $\forall \mathcal{F}, \bar{R}_T = o(T)$ .

Ex:  $A = \{1, \dots, K\}$  and  $\mathcal{F} = \underbrace{\left(\mathbb{P}([\alpha_1])\right)^K}_{\text{the case of finitely many atoms with bounded distributions}}$ , the set of all  $K$ -tuples of probability distributions over  $[\alpha_1]$

→ UCB does the job.

Counter-example:  $A = [\alpha_1]$  and  $\mathcal{F} = \underbrace{\left(\mathbb{P}([\alpha_1])\right)^{[\alpha_1]}}$

↑  
Illustrating that continuity is a minimal requirement.

all bandit problems  $(\mathbb{P}_x)_{x \in [\alpha_1]}$  with distributions  $\mathbb{P}_x$  having support  $[\alpha_1]$ .

Indeed: Consider  $(\delta_x)_{x \in [\alpha_1]}$  the bandit problem in which each arm  $x$  is associated with the Dirac mass on  $0$ .

Fix any strategy: it gets  $y_t = 0$  a.s. and

uses a sequence of (possibly) random choices  $I_t$ ,  $t \geq 1$

Since probability distributions can only have at most countably many atoms,

$$\mathcal{Y} = \{x \in [\alpha_1] : \exists t \mid \mathbb{P}\{I_t = x\} > 0 \text{ under } (\delta_x)_{x \in [\alpha_1]}\}$$

is countable. In particular,  $[\alpha_1] \setminus \mathcal{Y}$  is non-empty.

But the strategy behaves the same under the problem

$$(\tilde{\nu}_x^i)_{x \in [\alpha_1]} \quad \text{in which} \quad \begin{cases} \tilde{\nu}_x^i = \delta_0 & \forall x \neq x_0 \\ \tilde{\nu}_{x_0}^i = \delta_1 & \text{for one fixed } x_0 \in [\alpha_1] \setminus \mathcal{Y} \end{cases}$$

With probability 1, it thus never hits  $x_0$ .

Therefore,  $y_t = 0$  a.s. a.s. and  $\bar{R}_T = T - \mathbb{E}\left[\sum_{t=1}^T y_t\right] = T$ .

Actually, continuity is

sufficient for the regret to be controlled, as long as  $A$  is not too large.

Theorem: Let  $A$  be a <sup>cont</sup> metric space and let  $\mathcal{F}$  be the set of bandit problems  $(\mathbb{P}_x)_{x \in A}$

with: →  $\forall x$ ,  $\mathbb{P}_x$  is a distribution over  $[\alpha_1]$

→ a continuous mean-payoff function  $f: x \mapsto \mathbb{E}(\mathbb{P}_x)$

The regret can be controlled against  $\mathcal{F}^{\text{cont}}$  if and only if  $A$  is separable.

Corollary. Let  $\mathcal{F}^{\text{all}}$  be the family of all bandit models  $(\mathbb{P}_x)_{x \in A}$  with distributions  $\mathbb{P}_x$  over  $[0,1]$ . Then the regret against  $\mathcal{F}^{\text{all}}$  can be controlled if and only if  $A$  is at most countable.

Before we prove these facts, consider the following more concrete example, in which, by strengthening the regularity requirement on the mean-payoff function, we can even get rates.

Exercise :  
(Lipschitz bandits)

Let  $A = [0,1]$  and let  $\mathcal{F}^{\text{LP}}$  be the family of bandit models  $(\mathbb{P}_x)_{x \in [0,1]}$  with distributions  $\mathbb{P}_x$  over  $[0,1]$  and with mean-payoff functions that are Lipschitz.

Exhibit a strategy based on UCB + a sequence of discretizations of  $[0,1]$  into  $K$  bins (to be refined over time) such that:

Hint:

First, prove a performance bound by splitting  $[0,1]$  into  $[(i-1)/K, i/K]$  with  $i=1..K$

for a fixed  $K$ , where each bin  $[(i-1)/K, i/K]$  plays the role of an  $i$  in a bandit problem with finitely many arms.

Then discuss how to pick  $K$  over the time, as we do in the next proof.

$\forall x \in \mathcal{F}^{\text{LP}}$ ,

$$\bar{R}_T \leq (3L + 6\sqrt{8\ln T + 2})T^{2/3} + 2$$

where  $L$  is the Lipschitz constant of the mean-payoff function of  $x$ .

Proof of the Corollary :

We endow  $A$  with the discrete topology, ie, choose the distance  $d(x,y) = 1_{\{x \neq y\}}$ . Then:

1. All applications  $f: A \rightarrow \mathbb{R}$  are continuous
2.  $A$  is separable if and only if  $A$  is at most countable.

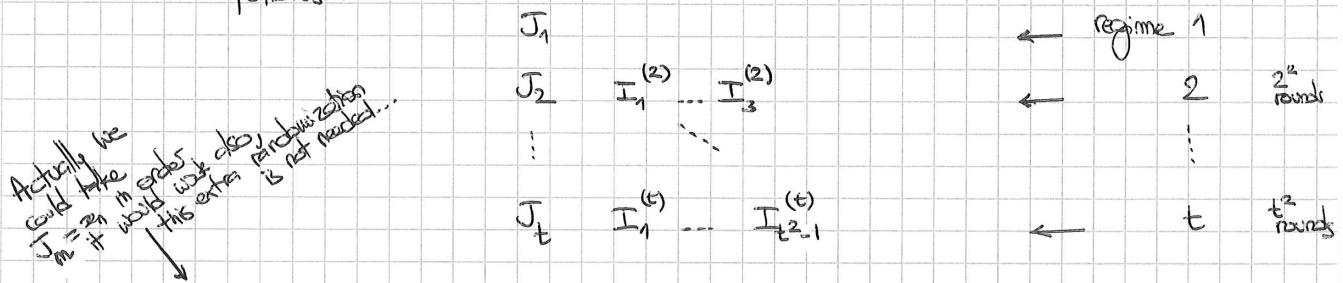
Proof of the Theorem:

It relies on the possibility or impossibility of uniform exploration of the arms.

1) If  $A$  is separable: let  $(x_n)_{n \in \mathbb{N}}$  be a collection of points in  $A$  that is dense

In particular, the probability distribution  $d = \sum_{n \geq 0} \frac{1}{2^{n+1}} \delta_{x_n}$  is such that  $d(V) > 0$  for all open sets  $V \subset A$ .

We pick elements  $J_1 J_2 I_1^{(2)} \dots J_t I_1^{(t)} \dots I_{t^2-1}^{(t)}$  as follows:



where the  $J_s$  are drawn at random according to  $d$  and the  $I_s^{(r)}$ ,  $1 \leq s \leq r^2-1$ , follow from the VCB strategy with arms  $J_1 \dots J_r$

$$\begin{aligned} \text{In regime } r: \quad & r^2 \max_{s \leq r} \mu_{J_s} - E \left[ \sum_{s=1}^{r^2} Y_{S_{r^2-1}+1} \right] \\ & \leq 1 + c \sqrt{r^3 \ln r} \end{aligned}$$

↑  
we have it first for  $E[Y_{J_1}]$  and then of course for  $E[Y_{J_r}]$  by tower rule

↑  
for  $J_r$   
well-chosen numerical constant  
distribution-free regret bound for VCB on  $r^2-1$  steps with  $r$  curves (we saw this bound as an exercise)

Let  $\varepsilon > 0$ , let  $\tilde{\tau}_\varepsilon$  the first (random) time when  $\mu_{J_r} = f(J_r) \geq \sup_{x \in A} f(x) - \varepsilon$

We have  $\tilde{\tau}_\varepsilon < \infty$  a.s. because:

- by continuity of  $f$ , there exists an open set  $V_\varepsilon$  with

$$\forall x \in V_\varepsilon, f(x) \geq \sup_A f - \varepsilon$$

- we have  $\tilde{r}_\Sigma \leq \inf \{ r \geq 1 : J_r \in V_\Sigma \} < +\infty$  a.s.

as this  $\uparrow$  random variable follows a geometric distribution with parameter  $d(V_\Sigma) > 0$ .

$$\text{For } r \geq \tilde{r}_\Sigma, \quad \max_{s \leq r} \mu_{J_s} + \varepsilon \geq \sup_A f$$

$$\text{So that } R_T = T \sup_A f - \mathbb{E}\left[\sum_{t=1}^T J_t\right]$$

$$\leq \sum_{r=1}^{\tilde{r}_\Sigma-1} r^2 + T\varepsilon + \sum_{r=\tilde{r}_\Sigma}^{r_T-1} r^2 \quad (1 + c\sqrt{r^3 \ln T})$$

lengths of  
regimes  $r \leq \tilde{r}_\Sigma-1$   
 $< +\infty$  a.s.

regime if  
may be  
incomplete

where  $r_T$  is such  
that time  $T$   
belongs to regime  $r_T$ :

we have  $r_T^3$  of  
order  $T$   
log  $r_T$  of order  $T^{1/3}$

and

$$\sum_{r \leq r_T} (1 + c\sqrt{r^3 \ln r})$$

$$\leq \sum_{r \leq r_T} (1 + cr^{3/2} \sqrt{\ln r}) = O(r_T^{5/2} \sqrt{\ln r_T}) \\ = O(T^{5/6} \sqrt{\ln T})$$

thus,

$$\limsup_{T \rightarrow +\infty} \frac{R_T}{T} \leq \varepsilon \quad \text{a.s.}$$

but since  $\bar{R}_T$  is a deterministic quantity  
and this is true  $\forall \varepsilon > 0$ , we have

$$\lim_{T \rightarrow +\infty} \frac{\bar{R}_T}{T} = 0 \quad \text{as requested.}$$

2) If  $A$  is not separable:

- \* We use the following characterization of separability (which relies on Zorn's lemma):

|| A metric space  $X$  is separable if and only if it contains no uncountable subset  $\mathcal{D}$  s.t.  $\rho = \inf \{ d(x_1, x_2) : x_1, x_2 \in \mathcal{D} \} > 0$ .

In particular, if  $A$  is not separable, there exist an uncountable subset  $\mathcal{D} \subset A$  and  $\rho > 0$  such that the balls  $B(a, \rho/2)$ , with  $a \in \mathcal{D}$ , are all disjoint.

$\hookrightarrow$  No probability distribution over  $A$  can give a positive mass to all these balls.

- \* We consider the bandit models  $\mathbb{J}^{(a)}$  inducing mean-payoff functions  $f^{(a)} : x \in A \mapsto (1 - \frac{d(x, a)}{\rho/2})^+$ ; in particular,  $\mathbb{J}_x^{(a)} = \delta_0$  for  $x \notin B(a, \rho/2)$ .  
 $\hookrightarrow f^{(a)}$  is indeed continuous.

We proceed as in the example showing the necessity of continuity when  $A = \mathbb{Q}$  and consider the bandit model  $(\delta_0)_{x \in A}$  as well as any strategy and the laws induced by the  $I_t$  under this model: let  $d_t$  be the law of  $I_t$  under  $(\delta_0)_{x \in A}$  and let  $d = \sum_{t \geq 1} \frac{1}{2^t} d_t$ .

There exists  $a \in A$  s.t.  $d(B(a, \rho/2)) = 0$ , that is, s.t.,  
 $\forall t \geq 1, \mathbb{P}(I_t \in B(a, \rho/2) \text{ under } (\delta_0)_{x \in A}) = 0$ .

The considered strategy is therefore such that the  $I_t$  have the same distribution under  $(\delta_0)_{x \in A}$  and  $\mathbb{J}^{(a)}$ . In particular,

$\mathbb{E}\left[\sum_{t=1}^T y_t\right] = 0$  in both cases, but in the latter case,  
 $\sup f^{(a)} = 1$ , so that  $\bar{R}_T = T$  against  $\mathbb{J}^{(a)}$ . The regret is not controlled against  $\mathbb{J}^{(a)} \in \mathbb{J}^{\text{cont}}$ .

Overview of the next steps: Fix a model  $\mathcal{D}$ , known to the decision-maker, ie a collection of probability distributions over  $\mathbb{R}$  with an expectation.

Assume that  $x_1, \dots, x_T$  are unknown but that the decision-maker knows  $\mathcal{D} = \{\mu_j\}_{j \in \mathcal{I}}$ .

What are the best bounds on  $\bar{R}_T = T\mu^* - \mathbb{E}\left[\sum_{t=1}^T Y_t\right]$ ?

We will show matching upper and lower bounds (with associated strategies):

$\bar{R}_T$  is at best of the order of  $\left(\sum_{a: \Delta_a > 0} \frac{\Delta_a}{\text{Kull}(x_a, \mu^*, \mathcal{D})}\right) \ln T$

where

$$\text{Kull}(x_a, \mu^*, \mathcal{D}) = \inf \left\{ \text{KL}(x_a^\dagger, x_a) : \begin{array}{l} x_a^\dagger \in \mathcal{D} \\ \mathbb{E}(x_a^\dagger) > \mu^* \end{array} \right\}$$

We will do so by

- proving a universal lower bound

Kullback  
Leibler  
divergence

expectation  
of  $x_a^\dagger$

- exhibiting a strategy, called KL-VG, to achieve the bound. ← if time permits (I'm not sure we will have time to do so...)

\* But \* before we do that, I guess that some reminder of basic and non-basic results about KL divergences would be needed!

The Kullback-Leibler divergence:definition and basic properties.

Definition (intrinsic): Let  $P, Q$  be two probability measures over  $(\Omega, \mathcal{F})$

$$KL(P||Q) = \begin{cases} +\infty & \text{if } P \text{ is not absolutely continuous w.r.t } Q \\ \int_{\Omega} \left( \frac{dP}{dQ} \ln \frac{dP}{dQ} \right) dQ = \int_{\Omega} \left( \ln \frac{dP}{dQ} \right) dP & \text{if } P \ll Q \end{cases}$$

Basic facts:

- Existence of the defining integral when  $P \ll Q$ : because  $\Psi: x \mapsto x \ln x$  is bounded from below on  $[0, +\infty)$
- $KL(P||Q) \geq 0$  and  $KL(P||Q) = 0 \iff P = Q$ .

It suffices to consider the case  $P \ll Q$ : because  $\Psi$  is strictly convex, Jensen's inequality indicates that

$$KL(P||Q) = \int_{\Omega} \Psi\left(\frac{dP}{dQ}\right) dQ \geq \Psi\left(\underbrace{\int \frac{dP}{dQ} dQ}_{=1}\right) = 0,$$

with equality if and only

if  $\frac{dP}{dQ}$  is  $Q$ -as constant, i.e.,  $P = Q$ .Exercise : A useful rewriting.

Prove the following result:

Assume  $P \ll Q$  and let  $\tau$  be any probability measure over  $(\Omega, \mathcal{F})$ such that  $P \ll \tau$  and  $Q \ll \tau$ . Denote  $f = \frac{dP}{d\tau}$  and  $g = \frac{dQ}{d\tau}$ .

$$\begin{aligned} \text{Then: } KL(P||Q) &= \int_{\Omega} \frac{f}{g} \ln\left(\frac{f}{g}\right) g d\tau \\ &= \int_{\Omega} \ln\left(\frac{f}{g}\right) f d\tau. \end{aligned}$$

Beware: with the usual measure-theoretic conventions, if  $x \neq 0$  and  $y=0$ , then  $\frac{x}{y} \neq \frac{y}{x}$  ↳ you therefore need to proceed with care!

Lemma (contraction of entropy; also known as data-processing inequality) :

Let  $P, Q$  be two probability measures over  $(\Omega, \mathcal{F})$

Let  $X: (\Omega, \mathcal{F}) \rightarrow (\Omega', \mathcal{F}')$  be any random variable

Denote by  $P^X$  and  $Q^X$  the laws of  $X$  under  $P$  and  $Q$ .

Then :

$$KL(P^X, Q^X) \leq KL(P, Q).$$

Proof: We may assume that  $P \ll Q$ , otherwise  $KL(P, Q) = +\infty$  and the inequality is true. We show that we then have

$$P^X \ll Q^X, \quad \text{with} \quad \frac{dP^X}{dQ^X} = E_Q \left[ \frac{dP}{dQ} \mid X = \cdot \right] \stackrel{\text{def.}}{=} \gamma$$

$$\text{i.e., } \gamma(x) = E_Q \left[ \frac{dP}{dQ} \mid X = x \right].$$

Indeed, for all  $B \in \mathcal{F}'$ :

tower rule

$$P^X(B) = P\{X \in B\} = \int_{\Omega} \mathbb{1}_B(x) \frac{dP}{dQ} dQ \stackrel{\downarrow}{=} \int_{\Omega} \mathbb{1}_B(x) E_Q \left[ \frac{dP}{dQ} \mid X = x \right] dQ$$

$$\stackrel{\text{def.}}{=} \int_{\Omega} \mathbb{1}_B(x) \gamma(x) dQ = \int_{\Omega} \mathbb{1}_B \gamma dQ^X.$$

Therefore,

$$\begin{aligned} KL(P^X, Q^X) &= \int_{\Omega} \gamma \ln \gamma dQ^X = \int_{\Omega} \gamma(x) \ln \gamma(x) dQ \\ &= \int_{\Omega} \left( E_Q \left[ \frac{dP}{dQ} \mid X \right] \ln E_Q \left[ \frac{dP}{dQ} \mid X \right] \right) dQ \quad \stackrel{\text{def.}}{\longrightarrow} \text{definition of } \gamma \\ &\leq \int_{\Omega} E_Q \left[ \frac{dP}{dQ} \ln \frac{dP}{dQ} \mid X \right] dQ \quad \stackrel{\text{conditional version of Jensen's inequality}}{\longleftarrow} \\ &\stackrel{\text{tower rule}}{\longrightarrow} = \int_{\Omega} \left( \frac{dP}{dQ} \ln \frac{dP}{dQ} \right) dQ = KL(P, Q) \end{aligned}$$

References: • The proof above is due to Ali and Silvey (1966), but it's far from being well-known!

- Typical proofs in the more recent literature:
  - either focus on the discrete case (Cover and Thomas, 2006)
  - or use the duality / variational formula for the KL (Massart, 2007; Boucheron, Lugosi, Massart, 2013)
- The joint convexity of KL, which we discuss below, is typically proved in a tedious way, relying on the rewriting of Exercise 1 and the joint convexity of  $(x,y) \in [0,+\infty)^2 \mapsto \left(\frac{x}{y} \ln \frac{x}{y}\right) y$

We may see it instead as a consequence of the data-processing inequality:

Corollary (joint convexity of KL): For all probability distributions  $P_1, P_2$  and  $Q_1, Q_2$  over the same measurable space  $(\Omega, \mathcal{F})$ , and all  $d \in (0,1)$ ,

$$\text{KL}\left((1-d)P_1 + dP_2, (1-d)Q_1 + dQ_2\right) \leq (1-d)\text{KL}(P_1, Q_1) + d\text{KL}(P_2, Q_2)$$

Proof: We augment  $(\Omega, \mathcal{F})$  into  $(\Omega \times \{1,2\}, \mathcal{F}')$  where

$$\mathcal{F}' = \mathcal{F} \otimes \{\emptyset, \{1\}, \{2\}, \{1,2\}\}$$

We define the random pair  $(X, J)$  by the projections

$$X: (\omega, j) \mapsto \omega \quad \text{and} \quad J: (\omega, j) \mapsto j$$

Let  $\bar{P}$  be a probability measure on  $(\Omega \times \{1,2\}, \mathcal{F}')$  such that

$$\begin{cases} J \sim 1 + \text{Ber}(d) \\ X | J=j \sim P_j \end{cases} \quad \begin{array}{l} \text{(and a similar definition for } Q \text{)} \\ \text{based on } Q_1, Q_2 \end{array}$$

that is,  $\forall j \in \{1,2\} \quad \forall A \in \mathcal{F} \quad \bar{P}(A \times \{j\}) = ((1-d)1_{\{j=1\}} + d1_{\{j=2\}}) P_j(A)$

$$\text{Now, } P^X = (1-d)P_1 + dP_2$$

$$Q^X = (1-d)Q_1 + dQ_2$$

and (as we prove below)  $\text{KL}(P^X, Q^X) = (1-d)\text{KL}(P_1, Q_1) + d\text{KL}(P_2, Q_2)$   
so that the result follows from the data-processing inequality.

Indeed: we may assume with no loss of generality, given  $d \in (0, 1)$ , that  
 $P_1 \ll Q_1$  and  $P_2 \ll Q_2$ , so that  $P \ll Q$  with

$$\frac{dP}{dQ}(w, j) = \mathbb{1}_{\{j=1\}} \frac{dP_1}{dQ_1}(w) + \mathbb{1}_{\{j=2\}} \frac{dP_2}{dQ_2}(w)$$

This entails that

$$\begin{aligned} \text{KL}(P, Q) &= \int_{\Omega \times \{1,2\}} \left( \frac{dP(w, j)}{dQ} \ln \frac{dP(w, j)}{dQ} \right) dQ(w, j) \\ &\quad \text{(we just use that for } f \geq 0 \text{ constant, } \int f dP = \int f \mathbb{1}_A dP) \\ &\quad + \text{if } f \text{ is the density of } P \text{ w.r.t. } Q \text{ then } f \text{ is measurable or not} \\ &= \int_{\Omega} \left( \frac{dP(w, 1)}{dQ} \ln \frac{dP(w, 1)}{dQ} \right) \mathbb{1}_{\{j=1\}}(w, j) dQ(w, j) \\ &\quad + \int_{\Omega} \left( \frac{dP(w, 2)}{dQ} \ln \frac{dP(w, 2)}{dQ} \right) \mathbb{1}_{\{j=2\}}(w, j) dQ(w, j) \\ &= \underbrace{\int_{\Omega} \left( \frac{dP_1}{dQ_1}(w) \ln \frac{dP_1}{dQ_1}(w) \right) dQ_1(w)}_{= \text{KL}(P_1, Q_1)} + \underbrace{\int_{\Omega} \dots}_{= \text{KL}(P_2, Q_2)} \end{aligned}$$

KL for product measures. ( $\leftrightarrow$  The independent case)

Proposition: Let  $(\Omega, \mathcal{F})$  and  $(\Omega', \mathcal{F}')$  be two measurable spaces,  
let  $P, Q$  be two probability measures over  $(\Omega, \mathcal{F})$   
 $P', Q'$  over  $(\Omega', \mathcal{F}')$

and denote by  $P \otimes P'$  and  $Q \otimes Q'$  the product distributions over  
 $(\Omega \times \Omega', \mathcal{F} \otimes \mathcal{F}')$ . Then:

$$\text{KL}(P \otimes P', Q \otimes Q') = \text{KL}(P, Q) + \text{KL}(P', Q')$$

Proof: We have  $P \triangleleft P' \Leftarrow Q \triangleleft Q' \Leftrightarrow [P \triangleleft Q \text{ and } P \triangleleft Q']$

so we can assume that all  $\ll$  statements hold, and then

$$\frac{d(P \otimes P')}{d(Q \otimes Q')} = \frac{dP}{dQ} \frac{dP'}{dQ'}$$

(this is a fundamental result in measure theory and one of the best characterizations of independence!).

Therefore,

$$KL(P \times P', Q \times Q') = \int_{\Omega} \left( \frac{dP}{dQ} \frac{dP'}{dQ'} \ln \left( \frac{dP}{dQ} \frac{dP'}{dQ'} \right) \right)$$

We see that  
if  $f$  &  $g$  are  
 $\geq$  a constant,  
then

$$\int (f+g) d\mu = \int f d\mu + \int g d\mu$$

Consequence (Girvier, Nérard, Stoltz, 2016) :

## Data-processing inequality with expectations of random variables

Corollary: Let  $P, Q$  be two probability measures over  $(\Omega, \mathcal{F})$

Let  $X: (\Omega, \mathcal{F}) \rightarrow ([\alpha_1], \mathcal{B}([\alpha_1]))$  be any  $[\alpha_1]$ -valued random variable.

Then, denoting by  $E_P[x]$  and  $E_Q[x]$  the respective expectations of  $X$

Under P and Q, we have:

$$E_P[x] \ln \frac{E_P[x]}{E_Q[x]} + (1 - E_P[x]) \ln \frac{1 - E_P[x]}{1 - E_Q[x]} = KL(Ber(E_P[x]), Ber(E_Q[x])) \leq KL(P, Q)$$

Proof: We denote by  $\mu$  the Lebesgue measure over  $[0,1]$  and augment the underlying measurable space into  $(\Omega \times [0,1], \mathcal{F} \otimes \mathcal{B}([0,1]))$ , over which we consider the product-distributions  $P \otimes \mu$  and  $Q \otimes \mu$ .

For any event  $E \in \mathcal{F} \otimes \mathcal{B}([0,1])$ , we have, by the data-processing inequality:

$$\begin{aligned}
 \text{KL}\left(\underbrace{(\mathbb{P} \otimes \eta)^{\perp E}}_{\text{Ber}(\mathbb{P} \otimes \eta(E))}, \underbrace{(\mathbb{Q} \otimes m)^{\perp E}}_{\text{Ber}(\mathbb{Q} \otimes m(E))}\right) &\leq \text{KL}(\mathbb{P} \otimes \eta, \mathbb{Q} \otimes m) \\
 &= \text{KL}(\mathbb{P}, \mathbb{Q}) + \text{KL}(\eta, m) \\
 &\stackrel{\uparrow}{=} \text{KL}(\mathbb{P}, \mathbb{Q})
 \end{aligned}$$

↑  
if product distributions

$$\text{Thus : } \text{KL}(\text{Ber}(\mathbb{P} \otimes \eta(E)), \text{Ber}(\mathbb{Q} \otimes m(E))) \leq \text{KL}(\mathbb{P}, \mathbb{Q})$$

The proof is concluded by picking  $E \in \mathcal{F} \otimes \mathcal{B}(\mathbb{R}^d)$  such that

$$\mathbb{P} \otimes \eta(E) = \mathbb{E}_{\mathbb{P}}[x] \quad \text{and} \quad \mathbb{Q} \otimes m(E) = \mathbb{E}_{\mathbb{Q}}[x]$$

Namely,  $E = \{(w, x) \in \mathbb{R} \times \mathbb{R}^d : x \leq x(w)\}$

By Tonelli's theorem :

$$\begin{aligned}
 \mathbb{P} \otimes \eta(E) &= \int_{\mathbb{R}} \left( \int_{\mathbb{R}^d} \mathbb{1}_{\{x \leq x(w)\}} d\eta(x) \right) d\mathbb{P}(w) \\
 &= \int_{\mathbb{R}} x(w) d\mathbb{P}(w) = \mathbb{E}_{\mathbb{P}}[x]
 \end{aligned}$$

and a similar equality for  $\mathbb{Q} \otimes m(E)$ .