

The MOSS strategy (Minimax Optimal Strategy in the Stochastic case of bandit problems)

Index policy relying on

$$U_a(t) = \hat{\mu}_a(t) + \sqrt{\frac{1}{2N_a(t)} \ln_+ \left(\frac{t}{K N_a(t)} \right)}$$

for $t \geq K$,

and where $\ln_+ = \max\{\ln, 0\}$

That is: For $t=1, \dots, K$: pull arm $A_t = t$

For $t \geq K+1$: pull arm $A_t \in \arg \max_{a=1 \dots K} U_a(t-1)$

Difference to UCB: we replace the exploration bonus

$$\sqrt{2 \ln t / N_a(t)}$$

by

$$\sqrt{\ln_+ \left(\frac{t}{K N_a(t)} \right) / (2 N_a(t))}$$

↳ no exploration after a
was pulled sufficiently
often (t/K times)

We prove a distribution-free bound:

Theorem: MOSS is such that $\sup_{\substack{J_1, \dots, J_K \\ \text{distributions} \\ \text{over } [0,1]}} \bar{R}_T \leq K-1 + 44\sqrt{KT}$

(the constant 44 can be improved)

but indeed
minimax optimal as
its name indicates!

$$= \int_0^{+\infty} \mathbb{P}\{Z_t^* \geq (\varepsilon u) N_t^*(t) \text{ and } N_t^*(t) \geq n_0\} du$$

where $Z_t^* = N_t^*(t) \left(\mu^* - \hat{\mu}_{N_t^*(t)} \right) = \sum_{s=1}^t \left(\mu^* - \frac{Y_s}{s} \right) \mathbb{1}_{\{N_s^*(t) = s^*\}}$
is a martingale,

and for all $x \in \mathbb{R}$, $S_{x,t} = e^{xZ_t^* - x^2/8 N_t^*(t)}$
is a supermartingale.

Thus by Markov-Chernoff, we continue the bounding as, for $x > 0$:

$$\leq \int_0^{+\infty} \mathbb{P}\{e^{xZ_t^* - x^2/8 N_t^*(t)} \geq \exp\left(N_t^*(t) \left(x(\varepsilon u) - \frac{x^2}{8}\right)\right) \text{ and } N_t^*(t) \geq n_0\} du$$

$$\leq \int_0^{+\infty} \sum_{l=n_0}^{+\infty} e^{-2l(\varepsilon u)^2} \mathbb{E}\left[S_{\frac{x}{4\varepsilon}t} \mathbb{1}_{\{N_t^*(t)=l\}}\right] du$$

we pick $x = 4(\varepsilon u)$
so that $x(\varepsilon u) - \frac{x^2}{8} = 2(\varepsilon u)^2$

⊗ independent of l , which
will be useful in other
proofs.

$$\leq e^{-2n_0\varepsilon^2} \int_0^{+\infty} e^{-2n_0u^2} \mathbb{E}\left[S_{\frac{x}{4\varepsilon}t} \mathbb{1}_{\{N_t^*(t) \geq n_0\}}\right] du$$

where $\mathbb{E}[S_{\frac{x}{4\varepsilon}t}] \leq 1$

we use $e^{-2l(\varepsilon u)^2} \leq e^{-2n_0(\varepsilon u)^2} \leq e^{-2n_0\varepsilon^2} e^{-2n_0u^2}$

All in all,

$$\mathbb{E}\left[(\mu^* - \hat{\mu}_{N_t^*(t)} - \varepsilon)^+ \mathbb{1}_{\{N_t^*(t) \geq n_0\}}\right]$$

$$\leq e^{-2n_0\varepsilon^2} \int_0^{+\infty} e^{-2n_0u^2} du = e^{-2n_0\varepsilon^2} \sqrt{\frac{\pi}{8n_0}}$$

where $\frac{\sqrt{\pi}}{\sqrt{8}} \leq 1$

integral of
a Gaussian density,
up to the normalization factor

Substituting the lemma in (*):

$$\mathbb{E}[\mu^* - U_{A^*}(t)] \leq \sqrt{\frac{K}{t}} + \sum_{l=0}^{+\infty} \underbrace{\frac{1}{\sqrt{x_{eH}}} e^{-2x_{eH} \frac{t}{Kx_e}}}_{\frac{1}{\sqrt{x_{eH}}} \exp\left(-2x_{eH} \frac{1}{2x_e} \ln\left(\frac{t}{Kx_e}\right)\right)} = \sqrt{\frac{K}{t}} \beta^{(l+1)/2} \exp\left(-\frac{l}{\beta} \ln \beta\right)$$

$\ln \beta^l = l \ln \beta$

where

$$\sum_{l=0}^{+\infty} \beta^{l(1/2 - 1/\beta)} \text{ is } < +\infty$$

as soon as $\beta \in (1, 2)$

$$= \sqrt{\frac{K}{t}} \beta^{(l+1)/2} \exp\left(-\frac{l}{\beta} \ln \beta\right)$$

$$= \sqrt{\frac{K}{t}} \beta^{1/2 + l(1/2 - 1/\beta)}$$

E.g., for $\beta = 3/2$,

$$\sum_{l=0}^{+\infty} \left(\frac{3}{2}\right)^{1/2 + l(1/2 - 2/3)} = \sqrt{\frac{3}{2}} \sum_{l=0}^{+\infty} \alpha^l = \frac{1}{1-\alpha} \sqrt{\frac{3}{2}} \leq 9$$

where $\alpha = \left(\frac{3}{2}\right)^{1/2 - 2/3}$

All in all: we obtain a $\sqrt{K/t} + 9\sqrt{K/t} = 20\sqrt{K/t}$ bound, as claimed.

Third step: $\sum_{t=K+1}^T \mathbb{E}[(U_{A_t}(t) - \mu_{A_t} - \sqrt{K/t})^+] \leq 4\sqrt{KT}$

$$= \sum_{t=K}^{T-1} \mathbb{E}[(U_{A_{t+1}}(t) - \mu_{A_{t+1}} - \sqrt{K/t})^+]$$

We decompose the expectations interest according to the $\{A_{t+1}=a\}$ and $\{N_a(t)=l\}$:

$$\sum_{t=K}^{T-1} \mathbb{E}[(U_{A_{t+1}}(t) - \mu_{A_{t+1}} - \sqrt{K/t})^+] = \sum_{a=1}^K \sum_{l=1}^{T-1} \sum_{t=K}^{T-1} \mathbb{E}[(U_a(t) - \mu_a - \sqrt{K/t})^+ \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_a(t)=l\}}]$$

We now use $(U_a(t) - \mu_a - \sqrt{K/t})^+ \leq (\hat{\mu}_a(t) - \mu_a - \sqrt{K/t})^+ + \begin{cases} 0 & \text{if } N_a(t) \geq \frac{t}{K} \\ \sqrt{\frac{2}{K} \ln\left(\frac{t}{K N_a(t)}\right)} & \text{if } N_a(t) < \frac{t}{K} \end{cases}$

also smaller than $\sqrt{\frac{1}{2 N_a(t)} \ln\left(\frac{T}{K N_a(t)}\right)}$ if $N_a(t) \leq \frac{t}{K}$

and get therefore the upper bound

$$\sum_{a=1}^K \sum_{l=1}^{T-1} \sum_{t=K}^{T-1} \mathbb{E}[(\hat{\mu}_a(t) - \mu_a - \sqrt{K/t})^+ \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_a(t)=l\}}]$$

$$+ \sum_{a=1}^K \sum_{l=1}^{T/K} \frac{1}{\sqrt{2l}} \ln\left(\frac{T}{Kl}\right) \mathbb{E}\left[\sum_{t=K}^{T-1} \mathbb{1}_{\{N_a(t)=l\}} \mathbb{1}_{\{A_{t+1}=a\}}\right]$$

We will repeatedly use that

$$\forall a, \ell_j \quad \sum_{t=K}^{T-1} \mathbb{1}_{\{N_t(t) = \ell_j\}} \mathbb{1}_{\{A_{t+1} = a\}} \leq 1 \quad (\text{ie, disjoint union})$$

cf. $N_t(t)$ increases by 1 whenever a is played

Also,
$$\sum_{\ell=1}^{T/K} \sqrt{\frac{1}{2\ell} \ln\left(\frac{T}{K\ell}\right)} \leq \int_0^{T/K} \sqrt{\frac{1}{2x} \ln\left(\frac{T}{Kx}\right)} dx$$

change of variable
 $u = T/(Kx)$

$$\leq \sqrt{\frac{T}{2K}} \int_1^{+\infty} u^{-3/2} \sqrt{\ln u} du$$

by the change of variable
 $u = e^{v^2}$

$$= \sqrt{\frac{T}{2K}} \int_0^{+\infty} 2v^2 e^{-v^2/2} dv = \sqrt{\pi} \sqrt{\frac{T}{K}}$$

Summarizing what we proved so far:

$$\sum_{t=K}^{T-1} \mathbb{E}[(U_{A_{t+1}}(t) - \mu_{A_{t+1}} - \sqrt{K_t})^+] \leq \sqrt{\pi} \sqrt{KT} + \sum_{a=1}^K \sum_{\ell=1}^{T/K} \sum_{t=K}^{T-1} \mathbb{E}[(\hat{\mu}_a(t) - \mu_a - \sqrt{K_t})^+ \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_t(t)=\ell\}}]$$

will show that $\leq \sqrt{\frac{\pi}{2}} \sqrt{\frac{T}{K}}$ for each a

We resort again to $Z_{a,t} = N_t(t) (\hat{\mu}_a(t) - \mu_a)$ martingale

and

$$S_{a,t}^{(a)} = e^{xZ_{a,t} - \frac{x^2}{2} N_t(t)}$$

supermartingale
where $x = 4(\sqrt{K_T} + u)$

For each a_j

$$\sum_{\ell=1}^{T/K} \sum_{t=K}^{T-1} \mathbb{E}[(\hat{\mu}_a(t) - \mu_a - \sqrt{K_t})^+ \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_t(t)=\ell\}}]$$

$$= \sum_{\ell=1}^{T/K} \sum_{t=K}^{T-1} \int_0^{+\infty} \mathbb{P}\left[xZ_{a,t} \geq N_t(t) \left(x(u + \sqrt{K_T}) - \frac{x^2}{2}\right) \mid A_{t+1}=a, N_t(t)=\ell\right] du$$

$$\stackrel{a=4(\sqrt{K_T}+u)}{\leq} \sum_{\ell=1}^{T/K} \sum_{t=K}^{T-1} \int_0^{+\infty} \underbrace{e^{-2\ell(u+\sqrt{K_T})^2}}_{\leq e^{-2\ell u^2}} e^{-2\ell K_T} \mathbb{E}\left[S_{a,t}^{(a)} \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_t(t)=\ell\}}\right] du$$

by Hoeffding-Chernoff

issue: this depends on t ...
but can be replaced in some sense by $S_{a,0}^{(a)} = 1$

the sum over t of these will be ≤ 1

But: remember Doob's maximal inequality for non-negative supermartingales:

$$\mathbb{P}\left\{\sup_{t \geq 0} S_{x,t}^{(a)} \geq c\right\} \leq \frac{\mathbb{E}[S_{x,0}^{(a)}]}{c} = \frac{1}{c}$$

→ see an alternative treatment on the next page

Then,

$$\sum_{l=1}^T \sum_{t=K}^{T-1} \mathbb{E}\left[\left(\hat{\mu}_a(t) - \mu_a - \sqrt{K_T}\right)^+ \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_h(t)=l\}}\right]$$

as before! →

$$= \sum_{l=1}^T \sum_{t=K}^{T-1} \int_0^{+\infty} \mathbb{P}\left\{S_{x,t}^{(a)} \geq e^{2\ell(u+\sqrt{K_T})^2} \text{ and } A_{t+1}=a \text{ and } N_h(t)=l\right\} du$$

$$\leq \sum_{l=1}^T \int_0^{+\infty} \sum_{t=K}^{T-1} \mathbb{P}\left\{\left(\sup_{s \geq 0} S_{x,s}^{(a)}\right) \geq e^{2\ell(u+\sqrt{K_T})^2} \text{ and } A_{t+1}=a \text{ and } N_h(t)=l\right\} du$$

$$\stackrel{\text{cf. disjoint union!}}{\leq} \sum_{l=1}^T \int_0^{+\infty} \mathbb{P}\left\{\left(\sup_{s \geq 0} S_{x,s}^{(a)}\right) \geq e^{2\ell(u+\sqrt{K_T})^2}\right\} du$$

Doob's maximal inequality

$$\sum_{l=1}^T \int_0^{+\infty} \underbrace{e^{-2\ell(u+\sqrt{K_T})^2}}_{\leq e^{-2\ell u^2} \times e^{-2\ell K_T}} du \leq \sum_{l=1}^T \frac{1}{\sqrt{\ell}} e^{-2\ell K_T}$$

and same treatment as in the lemma of the first part of the proof

This step is concluded by calculations:

$$\begin{aligned} \sum_{l=1}^T \frac{1}{\sqrt{\ell}} e^{-2\ell K_T} &\leq \int_0^T \frac{1}{\sqrt{x}} e^{-2x K_T} dx \\ &= \sqrt{\frac{T}{2K}} \int_0^{+\infty} \frac{e^{-u}}{\sqrt{u}} du = \sqrt{\frac{T}{2K}} \int_0^{+\infty} e^{-v^2} dv = \sqrt{\frac{\pi}{2}} \sqrt{\frac{T}{K}} \end{aligned}$$

Final bound is:

$$\sqrt{\pi} \sqrt{K_T} + K \sqrt{\frac{\pi}{2}} \sqrt{\frac{T}{K}} \leq 4 \sqrt{KT}$$

General conclusion: Final bound given by

$$\begin{aligned} K-1 + \left(\sum_{t=K+1}^T 2\sqrt{\frac{K}{t}}\right) + 4\sqrt{KT} &\leq K-1 + 4\sqrt{KT} + 2\sqrt{K} \int_0^T \sqrt{\frac{1}{t}} dt \\ &= K-1 + 44\sqrt{KT} \end{aligned}$$

Alternative treatment (credits to Enzo Miller) of the end of Step #3.

We were stuck at

$$\sum_{\ell=1}^T \sum_{k=K}^{T-1} \int_0^{+\infty} e^{-2\ell u^2} e^{-2\ell k/r} \mathbb{E} \left[S_{x,t}^{(a)} \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_k(t)=\ell\}} \right] du$$

$$= \sum_{\ell=1}^T \int_0^{+\infty} e^{-2\ell u^2} e^{-2\ell k/r} \mathbb{E} \left[\underbrace{\sum_{t=K}^{T-1} S_{x,t}^{(a)}}_{= S_{x,T_\ell}^{(a)}} \mathbb{1}_{\{A_{t+1}=a\}} \mathbb{1}_{\{N_k(t)=\ell\}} \right] du$$

where T_ℓ is given by:

$$T_\ell = \inf \{ t \geq 1 : A_{t+1} = a \text{ and } N_k(t) = \ell \}$$

We might get $\mathbb{E}[S_{x,T_\ell}^{(a)}] \leq \mathbb{E}[S_{x,0}^{(a)}] = 1$
 from the optional stopping theorem (« theoreme d'arrêt de Doob ») provided some verifications.

(In the first place: I am not so sure that T_ℓ is indeed a stopping time...)

Adversarial bandits.

(Rather stated in terms of losses than rewards!)

Setting:At each round $t=1,2,\dots$

1. The opponent and the decision-maker simultaneously choose $\ell_t = (\ell_{jt})_{j \in \{1, \dots, N\}}$ and $I_t \sim p_t$, where $p_t \in \mathcal{P}(\{1, \dots, N\})$
2. The opponent gets to see p_t and I_t ; the decision-maker only observes $\ell_{I_t t}$ (her own loss).

Regret:

$$R_T = \sum_{t=1}^T \ell_{I_t t} - \min_{j=1, \dots, N} \sum_{t=1}^T \ell_{jt}$$

vs. Pseudo-regret:

$$\bar{R}_T = \mathbb{E} \left[\sum_{t=1}^T \ell_{I_t t} \right] - \min_{j=1, \dots, N} \mathbb{E} \left[\sum_{t=1}^T \ell_{jt} \right]$$

↑
same definition as for stochastic bandits, up to the conversion of losses ℓ_{jt} into rewards $M - \ell_{jt}$ (for a well-chosen bound M)

↑
why $\mathbb{E}$?
cf. ℓ_{jt} are random variables, as they depend on the past, and in particular on $I_1, \dots, I_{t-1}$

We have $\bar{R}_T \leq \mathbb{E}[R_T]$.

We actually rather shoot for high-probability bounds on R_T , but studying $\bar{R}_T$ will be a good warm-up!



In these lecture notes, I'll take $N = K$ as the number of components

↳ we used N for individual sequences

↳ K stochastic bandits

and I alternatively took N and K in the next pgs...

(My bad...)

Adversarial bandits: bound on $\bar{R}_T$ via exponential weights.

Key: Estimators of the losses (the unseen and the seen ones):

$$\hat{\ell}_{jt} = \frac{\ell_{\mathbb{I}_t t}}{p_{jt}} \mathbb{1}_{\{\mathbb{I}_t = j\}} \quad \text{if } p_{jt} > 0 \quad \text{(which we will assume)}$$

auxiliary randomizations of opponent + decision-maker

They are (conditionally) unbiased: denoting by $\mathcal{F}_{t-1} = \sigma(U_1, \dots, U_{t-1}, \ell_1, \dots, \ell_{t-1}, p_1, \dots, p_{t-1}, \mathbb{I}_1, \dots, \mathbb{I}_{t-1})$

the total information available at the beginning of round t (of course, the decision-maker does not have that much information!), we have:

- ℓ_t and p_t are $\mathcal{F}_{t-1}$ -measurable; the only randomness comes from the random draw of $\mathbb{I}_t$ according to p_t
- $\hat{\ell}_{jt}$ can be rewritten $\hat{\ell}_{jt} = \frac{\ell_{jt}}{p_{jt}} \mathbb{1}_{\{\mathbb{I}_t = j\}}$

so that

$$\mathbb{E}[\hat{\ell}_{jt} | \mathcal{F}_{t-1}] = \frac{\ell_{jt}}{p_{jt}} \mathbb{E}[\mathbb{1}_{\{\mathbb{I}_t = j\}} | \mathcal{F}_{t-1}] = \frac{\ell_{jt}}{p_{jt}} p_{jt} = \ell_{jt}$$

since we assumed $p_{jt} > 0$

Algorithm: $p_1 = (1/N, \dots, 1/N)$ and for $t \geq 2$, $p_t = (p_{jt})_{j=1, \dots, N}$ is defined as

for a non-increasing sequence $(\eta_t)_{t \geq 2}$

$$p_{jt} = \exp\left(-\eta_t \sum_{s=1}^{t-1} \hat{\ell}_{js}\right) / \sum_{k=1}^N \exp\left(-\eta_t \sum_{s=1}^{t-1} \hat{\ell}_{ks}\right)$$

$\hookrightarrow$ ensures indeed that $p_{jt} > 0$.

the range $[0, M]$ is assumed to be known...

Theorem: The strategy above, tuned with $\eta_t = \frac{1}{M} \sqrt{\frac{\ln N}{Nt}}$, is such that:

for all opponents picking losses $\ell_{jt} \in [0, M]$,

$$\bar{R}_T = \mathbb{E}\left[\sum_{t=1}^T \ell_{\mathbb{I}_t t}\right] - \min_{i=1, \dots, N} \mathbb{E}\left[\sum_{t=1}^T \ell_{it}\right] \leq 2M \sqrt{TN \ln N}$$

The proof is based on the following lemma.

Lemma: The exponentially weighted average strategy on losses

$$\tilde{\ell}_{jt} \in [0, +\infty[, \quad \text{i.e.,}$$

$$\text{with } \eta_t \downarrow, \quad \tilde{p}_{jt} = \frac{\exp(-\eta \sum_{s=1}^{t-1} \tilde{\ell}_{js})}{\sum_{k=1}^N \exp(-\eta \sum_{s=1}^{t-1} \tilde{\ell}_{ks})},$$

is such that

$$\sum_{t=1}^T \sum_{j=1}^N \tilde{p}_{jt} \tilde{\ell}_{jt} - \min_{i=1 \dots N} \sum_{t=1}^T \tilde{\ell}_{it} \leq \frac{\ln N}{\eta_T} + \sum_{t=1}^T \frac{\eta_t}{2} \sum_j \tilde{p}_{jt} \tilde{\ell}_{jt}^2.$$

Proof: We saw earlier in this series of lectures that the EWA strategy (with $\eta_t \downarrow$) is such that

$$\forall \tilde{\ell}_{jt} \in \mathbb{R}, \quad \sum_{j=1}^N \tilde{p}_{jt} \tilde{\ell}_{jt} - \min_{i=1 \dots N} \sum_{t=1}^T \tilde{\ell}_{it} \leq \frac{\ln N}{\eta_T} + \sum_{t=1}^T \tilde{\mathcal{S}}_t$$

$$\text{where } \tilde{\mathcal{S}}_t = \sum_{j=1}^N \tilde{p}_{jt} \tilde{\ell}_{jt} + \frac{1}{\eta_t} \ln \sum_{j=1}^N \tilde{p}_{jt} e^{-\eta_t \tilde{\ell}_{jt}}$$

$$\text{We use here } e^{-x} \leq 1 - x + \frac{x^2}{2} \quad \forall x \geq 0$$

$$\begin{aligned} \text{so that } \ln \sum_j \tilde{p}_{jt} e^{-\eta_t \tilde{\ell}_{jt}} &\leq \ln \left(1 - \eta_t \sum_j \tilde{p}_{jt} \tilde{\ell}_{jt} + \frac{\eta_t^2}{2} \sum_j \tilde{p}_{jt} \tilde{\ell}_{jt}^2 \right) \\ &\stackrel{\substack{\ln(1+x) \leq x \\ \forall x > -1}}{\leq} - \eta_t \sum_j \tilde{p}_{jt} \tilde{\ell}_{jt} + \frac{\eta_t^2}{2} \sum_j \tilde{p}_{jt} \tilde{\ell}_{jt}^2 \end{aligned}$$

hence the stated bound.

Proof (of the Theorem): We have no control on how large the $\tilde{\ell}_{jt}$ can be, and they could be very large! So, we would not be ready to apply any bound with a remainder $M_T \ln N$ term, where M_T is such that $\tilde{\ell}_{jt} \in [0, M_T]$ $\forall j, t$... as this M_T could be even super-linear. That's why we go back to the beginning of the

proof for the fully adaptive algorithm The η_t can be picked as $\ln N / \sum_{s=1}^{t-1} \delta_s$ or in terms of the upper bounds on the δ_s (we choose the latter version for the sake of concreteness).
 The lemma yields for the $\hat{L}_{jt}$: $\hookrightarrow$ see belows!

$$\sum_{t,j} p_{jt} \hat{L}_{jt} - \min_{i=1,\dots,N} \sum_{t=1}^T \hat{L}_{it} \leq \frac{\ln N}{\eta_T} + \sum_{t=1}^T \frac{\eta_t}{2} \sum_j p_{jt} \hat{L}_{jt}^2$$

by definition of the $\hat{L}_{jt}$ similar treatment:

$$= \sum_j L_{jT} \frac{p_{jt}}{p_{jt}} \mathbb{1}_{jT=jT} = L_{jT}$$

$$= \sum_j L_{jT}^2 \frac{1}{p_{jt}} \mathbb{1}_{jT=jT} \leq M^2 \sum_j \mathbb{1}_{jT=jT} / p_{jt}$$

To simplify even further the choice of the η_t , we first take $\mathbb{E}$ of both sides:

$$\mathbb{E} \left[\sum_t L_{jT} \right] - \mathbb{E} \left[\min_i \sum_t \hat{L}_{it} \right] \leq \frac{\ln N}{\eta_T} + \frac{M^2}{2} \sum_{t=1}^T \eta_t \sum_{j=1}^N \underbrace{\mathbb{E} \left[\frac{\mathbb{1}_{jT=jT}}{p_{jt}} \right]}_{=1}$$

$\leq \min_i \sum_{t=1}^T \mathbb{E}[\hat{L}_{it}] = \mathbb{E}[\hat{L}_{it}]$ by the tower rule and the fact that $\hat{L}_{it}$ is conditionally unbiased.

Thus,

$$\bar{R}_T = \mathbb{E} \left[\sum_{t=1}^T L_{jT} \right] - \min_{i=1,\dots,N} \mathbb{E} \left[\sum_{t=1}^T \hat{L}_{it} \right] \leq \frac{\ln N}{\eta_T} + \frac{M^2 N}{2} \sum_{t=1}^T \eta_t$$

The only adaptation to be made is w.r.t T (as M is assumed to be known)

The optimal constant η would be s.t. $\ln N / \eta = \frac{M^2 N}{2} T \eta$, that is,

$$\eta \text{ proportional to } \frac{1}{M} \sqrt{\frac{\ln N}{NT}}$$

$$\hookrightarrow \text{try } \eta_t = \frac{\gamma}{M} \sqrt{\frac{\ln N}{Nt}} \quad \text{where } \gamma \text{ is to be optimized.}$$

The bound is
$$M \sqrt{N \ln N} \left(\frac{\sqrt{T}}{\gamma} + \gamma \sum_{t=1}^T \frac{1}{\sqrt{t}} \right)$$

$$\leq \int_0^T \frac{1}{\sqrt{t}} dt \leq 2\sqrt{T}$$

$$\leq \sqrt{T} \left(\frac{1}{\gamma} + \gamma \right) = 2\sqrt{T} \quad \text{for } \gamma=1$$

Remarks / insights:

0 to apply
safely
 $e^{-x} \leq 1 - x + \frac{x^2}{2}$

M to bound $\ell_{i,t}^2 \leq M^2$
and to pick η_t

→ The range $[0, M]$ needs to be known; it could be a general range $[m, M]$, in which case, we would translate all losses by $-m$,
eg, consider $\hat{\ell}_{i,t} = \frac{\ell_{i,t} - m}{M - m} \mathbb{1}_{\{I_t = j\}}$

and $\eta_t = \frac{1}{M-m} \sqrt{\frac{\ln N}{Nt}}$

to get $R_T \leq 2(M-m) \sqrt{T N \ln N}$

↳ Adaptation to the range in adversarial bandits is actually a difficult issue (I would need to think more about it...)

→ In the next page you'll get a brief view of how to get high probability bounds on the true regret, of the same order of magnitude:

wp $1-\delta$, $R_T \leq (M-m) \sqrt{T N \ln(N/\delta)}$

→ There exists an algorithm (INF) st. R_T is controlled (in $\mathbb{E}$ or with high probability) against individual sequences $\ell_{i,t}$ [ie: sequences fixed in advance, not chosen by an opponent who reacts]

↳ « oblivious » opponent

The $\ell_{i,t}$ are not random variables, in this case.